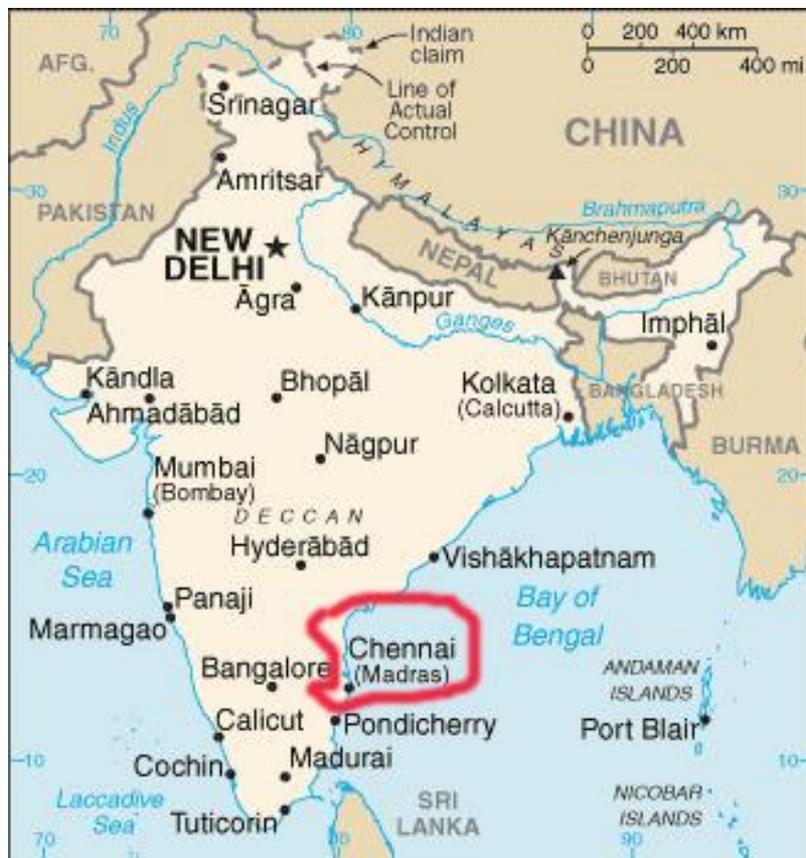


Attacking KLEIN

J.-Ph. Aumasson, María Naya-Plasencia,
Markku-Juhani O. Saarinen

Presented at INDOCRYPT 2011 (December)



KLEIN: RFIDSec 2011 (June)

rfid-cusp.org/rfidsec/program.php

10:10–10:35 AM Coffee and Tea Break

Technical Session 1: On-Tag Cryptography
(Session chair: Farinaz Koushanfar)

KLEIN: A New Family of Lightweight Block Ciphers

Authors: Zheng Gong, Svetla Nikova and Yee Wei Law

Affiliations: School of Computer Science, South China Normal University, China and Faculty of EWI, University of Twente, The Netherlands and Department of EEE, The University of Melbourne, Australia

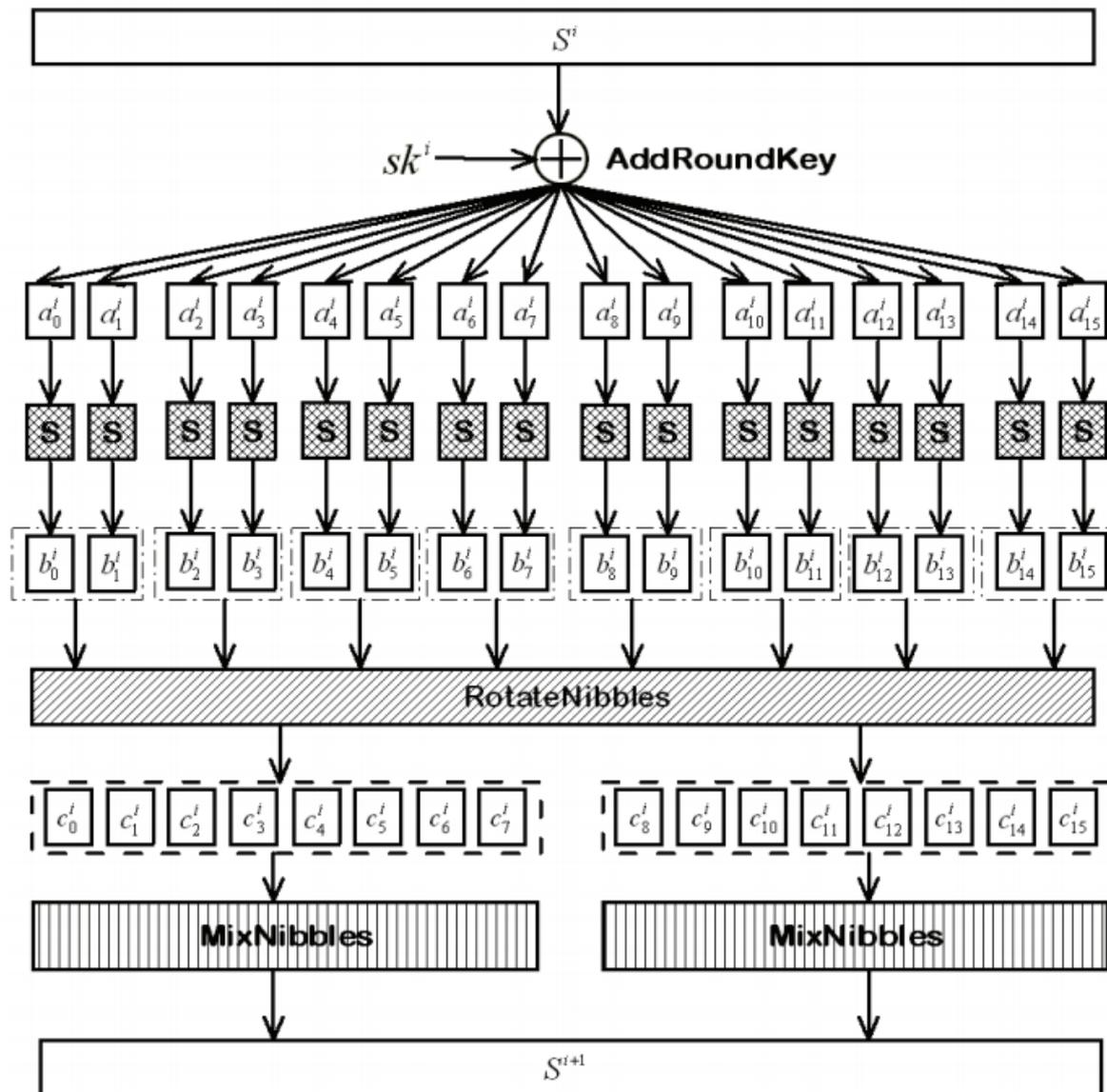
Low-RAM on small MCUs

~2000 GE on 180 nm ASIC

Reasonable speed

Simple and elegant design

=> How secure is it?



64-bit block

64-bit key -> 12 rounds

80-bit key -> 16 rounds

96-bit key -> 20 rounds

Interaction between:

Nibble-oriented S-box

Byte-oriented MixColumn

Observation 1

$X \in \{1,2,\dots,7\}$ [i.e. nibble with null MSB]

$\text{MixColumn}(m) \oplus \text{MixColumn}(m \oplus 0000000X)$

$= 0Y0Y0Y0Y$

Where Y is a wildcard for nibbles $\in \{1,2,\dots,F\}$

Observation 2

X **wildcard** for nibbles $\in \{1,2,\dots,7\}$

MixColumn(m) \oplus MixColumn(m \oplus 0X0X0X0X)

= 0Y0Y0Y0Y

Where Y is a wildcard for nibbles $\in \{0,1,\dots,F\}$

Observation 3

X wildcard for nibbles $\in \{8,9,\dots,F\}$

$\text{MixColumn}(m) \oplus \text{MixColumn}(m \oplus \text{0X0X0X0X})$

$= \text{0Y0Y0Y0Y}$

Where Y is a wildcard for nibbles $\in \{0,1,\dots,F\}$

Observation 4

$$X \in \{1,2,\dots,F\}$$

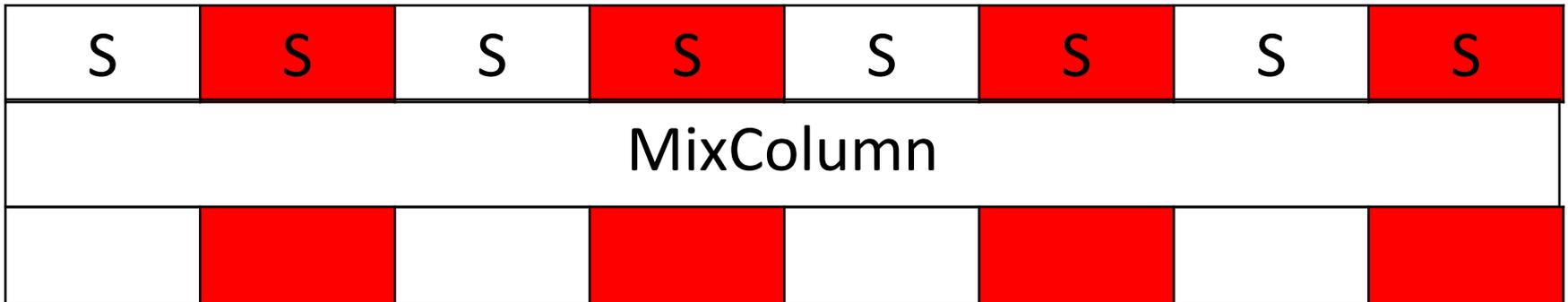
$$\text{Prob}[\text{Sbox}(m) \oplus \text{Sbox}(m \oplus x)] \in \{1,2,\dots,8\}$$

$$= 7/15 \sim 2^{-1.1}$$

$$\text{If } X \in \{B,E\}, \text{ Prob} = 3/4$$

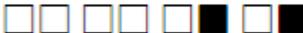
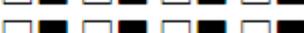
Key idea

Ensure that differences entering MixColumn are all either in $\{0,1,\dots,7\}$ or in $\{8,9,\dots,F\}$, so that its output only has differences in lower nibbles



1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
2	SubNibbles RotateNibbles MixNibbles			$p_2 \approx 2^{-4.40}$	$2^{-4.82}$
3	SubNibbles RotateNibbles MixNibbles			$p_3 \approx 2^{-5.82}$	$2^{-10.64}$
4	SubNibbles RotateNibbles MixNibbles			$p_4 \approx 2^{-5.82}$	$2^{-16.45}$
5	SubNibbles RotateNibbles MixNibbles			$p_5 \approx 2^{-5.82}$	$2^{-22.27}$
6	SubNibbles RotateNibbles MixNibbles			$p_6 \approx 2^{-5.82}$	$2^{-28.08}$
7	SubNibbles RotateNibbles MixNibbles			$p_7 \approx 2^{-5.82}$	$2^{-33.90}$

Black squares = actives nibbles

1	SubNibbles RotateNibbles MixNibbles	  	  	$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
2	SubNibbles RotateNibbles MixNibbles	  	  	$p_2 \approx 2^{-4.40}$	$2^{-4.82}$
3	SubNibbles RotateNibbles MixNibbles	  	  	$p_3 \approx 2^{-5.82}$	$2^{-10.64}$
4	SubNibbles RotateNibbles MixNibbles	  	  	$p_4 \approx 2^{-5.82}$	$2^{-16.45}$
5	SubNibbles RotateNibbles MixNibbles	  	  	$p_5 \approx 2^{-5.82}$	$2^{-22.27}$
6	SubNibbles RotateNibbles MixNibbles	  	  	$p_6 \approx 2^{-5.82}$	$2^{-28.08}$
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Distinguisher on 6 rounds in $\sim 2^{28}$

1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
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Distinguisher on 6 rounds in $\sim 2^{28}$
 Transform to key-recovery on 7 rounds

1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
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7	SubNibbles RotateNibbles MixNibbles			$p_7 \approx 2^{-5.82}$	$2^{-33.90}$

Transform to key-recovery on 7 rounds

Message modification $\Rightarrow 2^{23}$ for an extra pair

1	SubNibbles RotateNibbles MixNibbles	  	  	$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
2	SubNibbles RotateNibbles MixNibbles	  	  	$p_2 \approx 2^{-4.40}$	$2^{-4.82}$
3	SubNibbles RotateNibbles MixNibbles	  	  	$p_3 \approx 2^{-5.82}$	$2^{-10.64}$
4	SubNibbles RotateNibbles MixNibbles	  	  	$p_4 \approx 2^{-5.82}$	$2^{-16.45}$
5	SubNibbles RotateNibbles MixNibbles	  	  	$p_5 \approx 2^{-5.82}$	$2^{-22.27}$
6	SubNibbles RotateNibbles MixNibbles	  	  	$p_6 \approx 2^{-5.82}$	$2^{-28.08}$
7	SubNibbles RotateNibbles MixNibbles	  	  	$p_7 \approx 2^{-5.82}$	$2^{-33.90}$

Message modification $\Rightarrow 2^{23}$ for an extra pair
 Find 6 pairs in $<2^{29}$, enough to find 32b subkey

1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
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6	SubNibbles RotateNibbles MixNibbles			$p_6 \approx 2^{-5.82}$	$2^{-28.08}$
7	SubNibbles RotateNibbles MixNibbles			$p_7 \approx 2^{-5.82}$	$2^{-33.90}$

Find 6 pairs in $<2^{29}$, enough to find 32b subkey
Repeat for the other 32b half

1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
2	SubNibbles RotateNibbles MixNibbles			$p_2 \approx 2^{-4.40}$	$2^{-4.82}$
3	SubNibbles RotateNibbles MixNibbles			$p_3 \approx 2^{-5.82}$	$2^{-10.64}$
4	SubNibbles RotateNibbles MixNibbles			$p_4 \approx 2^{-5.82}$	$2^{-16.45}$
5	SubNibbles RotateNibbles MixNibbles			$p_5 \approx 2^{-5.82}$	$2^{-22.27}$
6	SubNibbles RotateNibbles MixNibbles			$p_6 \approx 2^{-5.82}$	$2^{-28.08}$
7	SubNibbles RotateNibbles MixNibbles			$p_7 \approx 2^{-5.82}$	$2^{-33.90}$

7 round distinguisher? $33.90 > 32...$

Filter by checking ability to message-modify a pair

1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
2	SubNibbles RotateNibbles MixNibbles			$p_2 \approx 2^{-4.40}$	$2^{-4.82}$
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7 round distinguisher? $33.90 > 32...$

Filter by checking ability to message-modify a pair

1	SubNibbles RotateNibbles MixNibbles	 	 	$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
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7 round distinguisher? $33.90 > 32...$

If trail followed, an extra pair costs $\sim 2^{28}$

1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
2	SubNibbles RotateNibbles MixNibbles			$p_2 \approx 2^{-4.40}$	$2^{-4.82}$
3	SubNibbles RotateNibbles MixNibbles			$p_3 \approx 2^{-5.82}$	$2^{-10.64}$
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6	SubNibbles RotateNibbles MixNibbles			$p_6 \approx 2^{-5.82}$	$2^{-28.08}$
7	SubNibbles RotateNibbles MixNibbles			$p_7 \approx 2^{-5.82}$	$2^{-33.90}$

A good pair can be identified in $\sim 2^{34}$ encryptions
 Find 32b half subkey that gives inactive higher nibbles

1	SubNibbles RotateNibbles MixNibbles			$p_1 \approx 2^{-0.42}$	$2^{-0.42}$
2	SubNibbles RotateNibbles MixNibbles			$p_2 \approx 2^{-4.40}$	$2^{-4.82}$
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7	SubNibbles RotateNibbles MixNibbles			$p_7 \approx 2^{-5.82}$	$2^{-33.90}$

Repeat for the other half of the last subkey
=> Recover the full 64b key in $<2^{35}$ encryptions

8-round key recovery

$<2^{35}$ for 64b key (12 rounds in total)

$<2^{51}$ for 80b key (16 rounds in total)

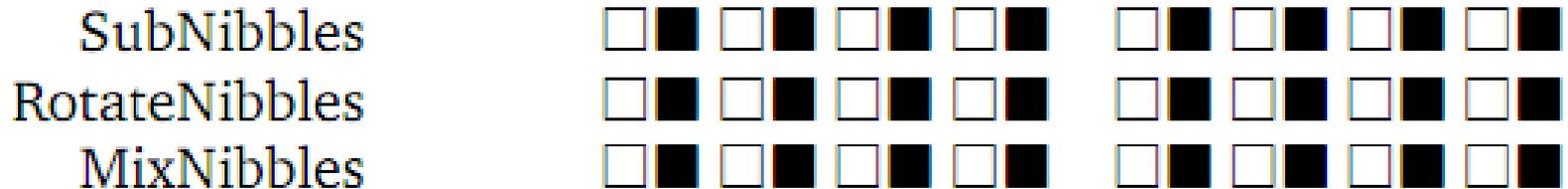
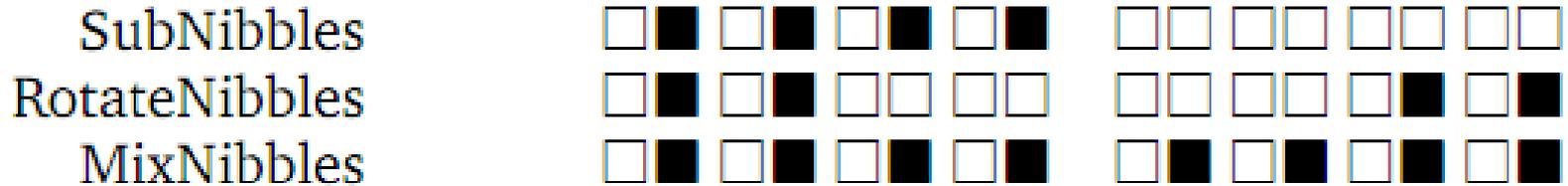
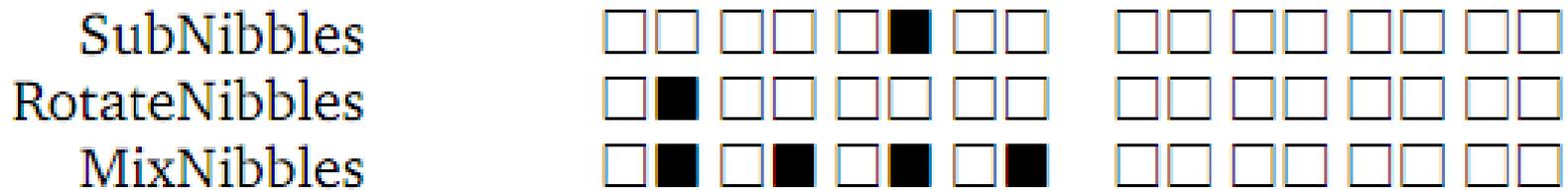
$<2^{67}$ for 96b key (20 rounds in total)

```
$ ./attack 8
test vector ok
soundness ok
Pair found in 2^28.21: fb5248c1a424ca3e
Pair found in 2^26.43: 00b848c1a424882f
Pair found in 2^28.54: 180b48c1a4245a09
Pair found in 2^26.78: 1ee948c1a4246b1d
Pair found in 2^25.81: 226848c1a424362e
Pair found in 2^27.56: 2e3548c1a424f161
Subkey lower nibbles recovered:
d42c
d515
Actual subkey lower nibbles:
d42c d515
1344 seconds elapsed
```

Extend to 9 rounds?

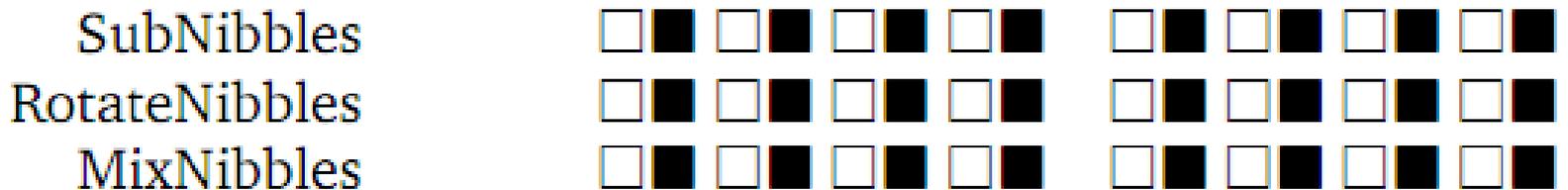
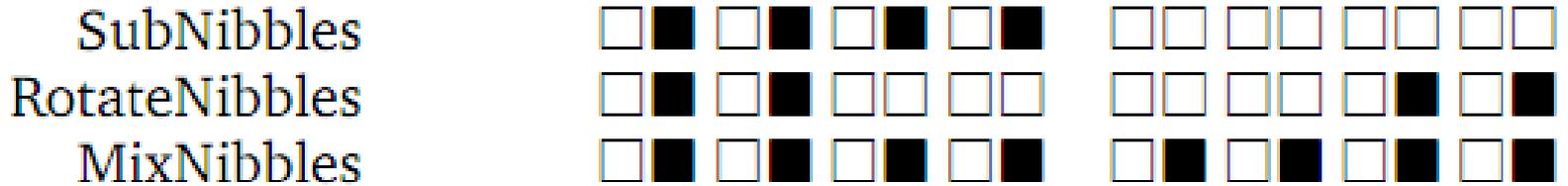
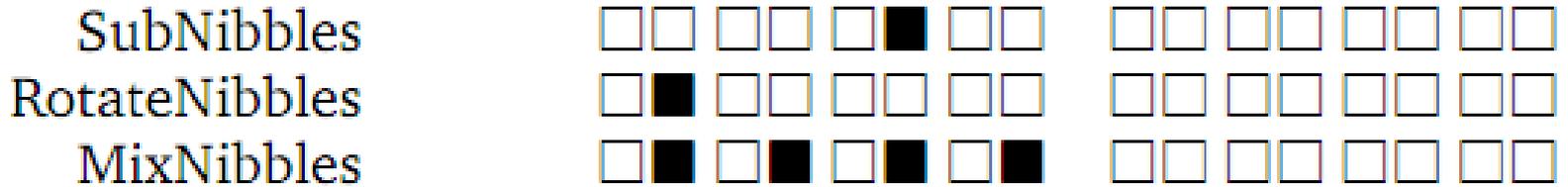
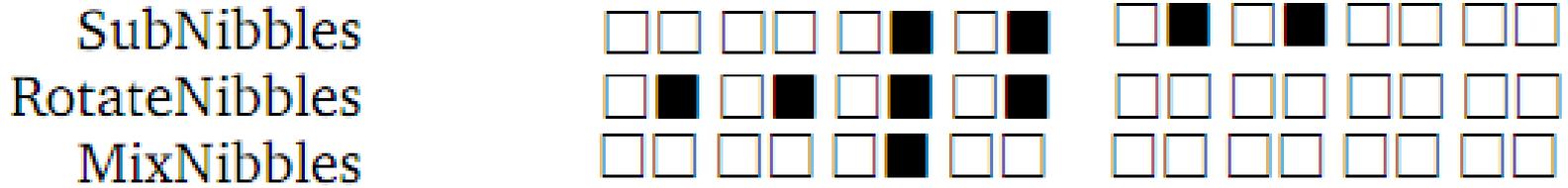
(work in progress with Y. Sasaki)

Prepend a round here

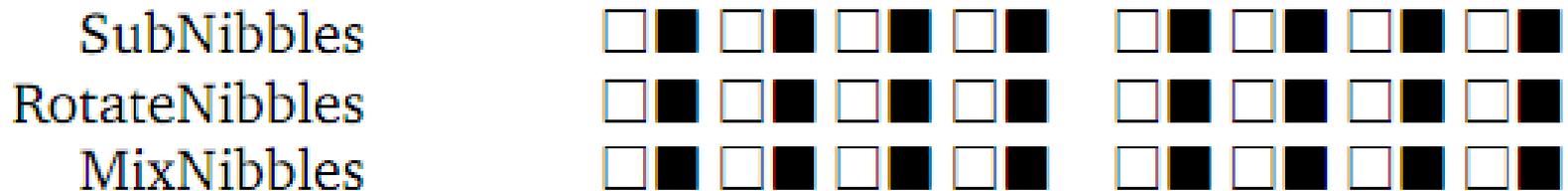
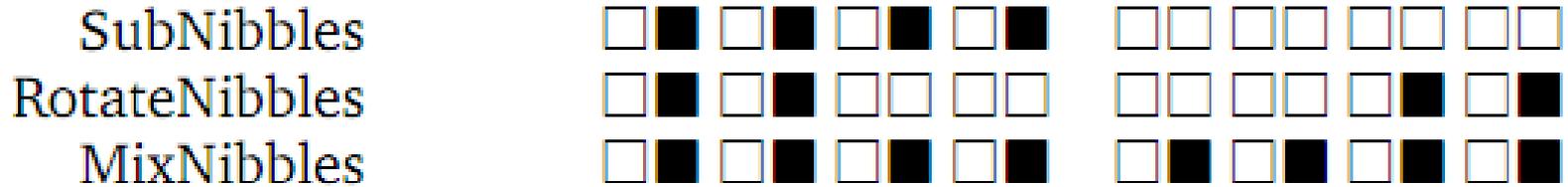
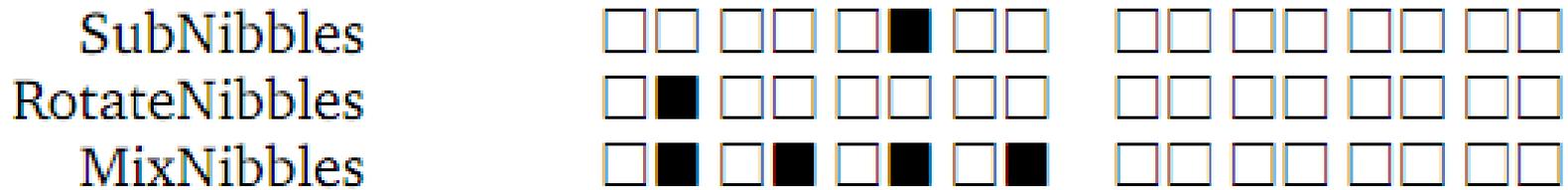
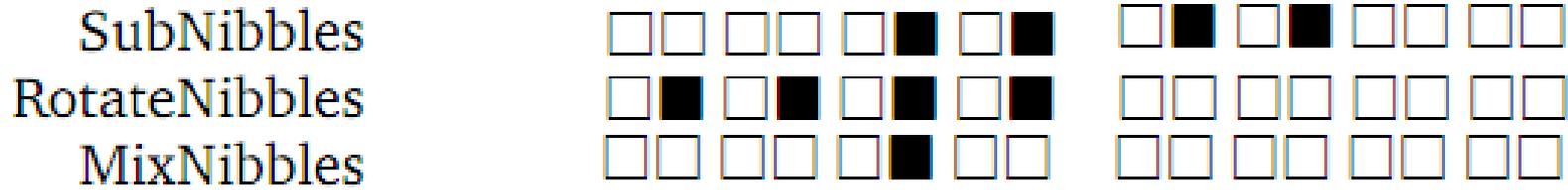


Guess 16 subkey bits of first subkey

Success \Leftrightarrow good post-SubNibbles difference

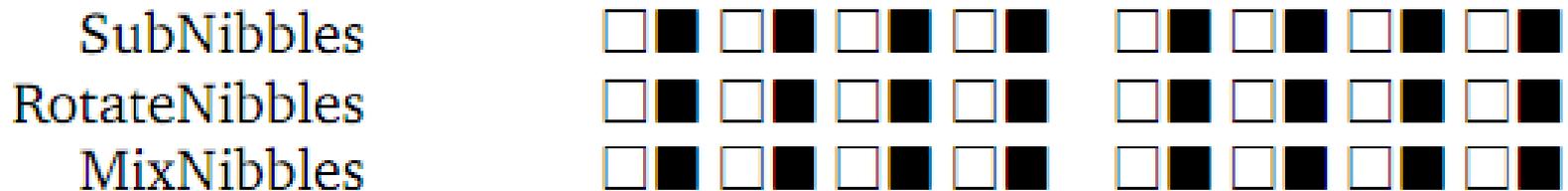
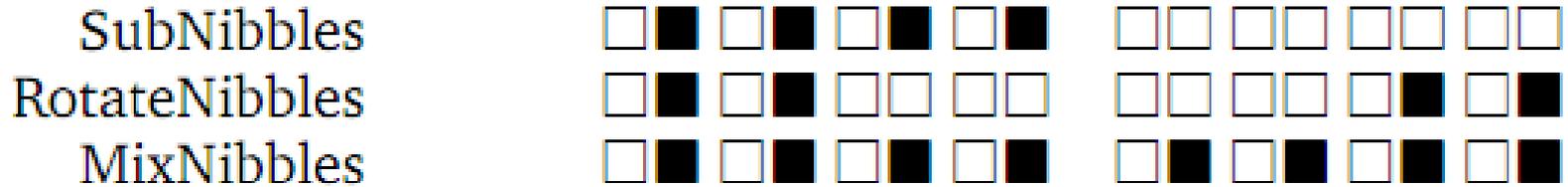
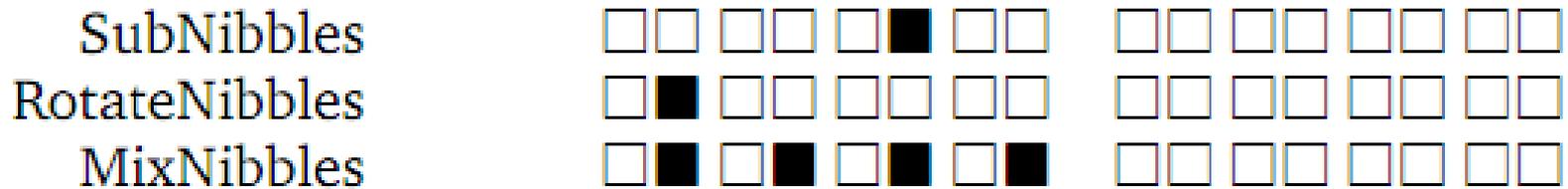
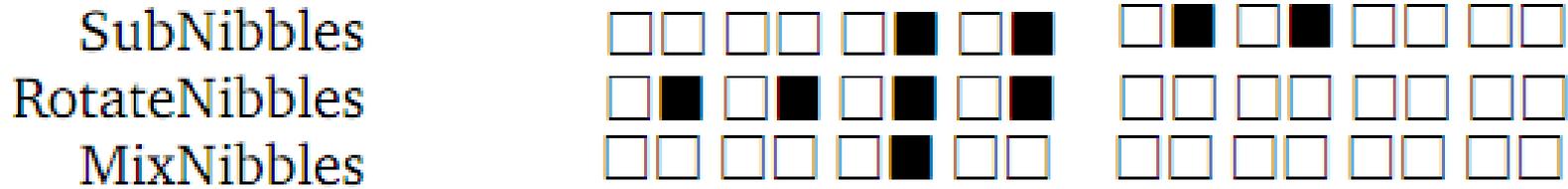


Correct guess + correct pair => 1st round passed
 Use structure to reduce data cplexity



After 1st round, same attack as on 8 rounds, without message modification (filtering still possible)

To be verified...



8-round attacks claimed at INSCRYPT 2011

Similar first-round trick

Unverified probabilities

INSCRYPT 2011

In Cooperation with IACR

Nov. 30 - Dec. 3, 2011, Beijing

	Feng	and F-FCSR-H v3
48	Nan Li, Yi Mu and Willy Susilo	Efficient Self-Certified Signatures with Batch Verification
49	Jiang Zhang, Xiang Xie, Rui Zhang and Zhenfeng Zhang	A Generic Construction from Selective-IBE to Public-Key Encryption with Non-interactive Opening
53	Xusheng Zhang, Shan Chen and Dongdai Lin	Fast Tate Pairing Computation on Twisted Jacobi Intersections Curves
54	Shao-Zhen Chen and Yi-Bin Dai	Weak-Key Class of MISTY1 for Related-Key Differential Attack
55	Xiaoli Yu, Wenling Wu, Yanjun Li and Lei Zhang	Cryptanalysis of Reduced-Round KLEIN Block Cipher

TODO

- Analyze/implement 9-round attack
- What about hashing modes (DM etc.)?
- 10+ rounds can probably be attacked in $<2^{64}$

